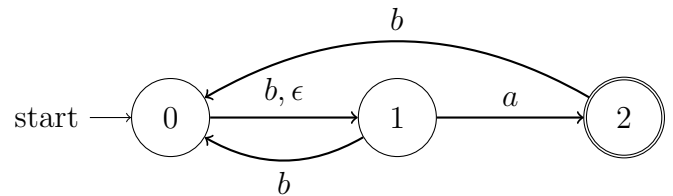


Due Monday, September 18.

- Convert the following NFA to a DFA. Use the method we discussed in class, where the states of the DFA correspond to subsets of the states of the original NFA. Hint: *After removing states in the DFA that you can never reach, you should only need a small number of states, one of which corresponds to the empty set.*



- Let  $\Sigma = \{0, 1\}$ . Write a one-sentence description the languages defined by the following regular expressions. For example:  $\Sigma^*1$  would be any binary string that ends with a 1.

(a)  $(\Sigma\Sigma)^*$ .

(b)  $\Sigma^*01\Sigma^*$ .

(c)  $(0\Sigma^*0)|(1\Sigma^*1)$ .

- Find a regular expression that matches each of the following languages. In all cases, the alphabet is  $\Sigma = \{0, 1\}$ .

(a)  $\{w \in \Sigma^* : w \text{ contains at least three 1's.}\}$

(b)  $\{w \in \Sigma^* : w \text{ contains at least two 1's and exactly one 0.}\}$

4. Let  $\Sigma$  be the regular English alphabet  $\{a, b, c, \dots, z\}$ . Write a regular expression that matches all strings that contain at least two vowels (i.e.,  $a, e, i, o, u$ ).
5. Prove that if  $L \subset \Sigma^*$  is a regular language, then complement  $\Sigma^* \setminus L$  is also a regular language. Hint: *If there is a DFA  $M = (Q, \Sigma, \delta, q, F)$  that recognizes  $L$ , describe a different DFA that recognizes the complement.*
6. Let  $L = \{w \in \Sigma^* : \text{the length of } w \text{ is a power of } 2\}$ . Use the pumping lemma to prove that  $L$  is not a regular language.