Introduction to Compiler Design Lecture 1 Chapters 1 and 2

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The Stages of Compilation

- Lexical Analysis
- Syntactic Analysis
- Semantic Analysis
- Intermediate Code Generation
- Optimization
- Machine Code Generation

2 Assignment

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2 Assignment

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The Stages of Compilation

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2 Assignment

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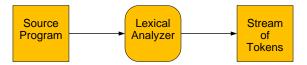
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Definition (Token)

A token is a smallest meaningful group symbols.

Definition (Lexical analyzer)

A lexical analyzer, also called a lexer or a scanner, receives a stream of characters from the source program and groups them into tokens.



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Example (Lexical Analysis)

• What are the tokens in the following program?

```
int main()
{
    float a = 123.4;
    return 0;
}
```

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- Each token has a type and a value.
- For example,
 - The variable count has type id and value "count".
 - The number 123 has type num and value "123".
 - The keyword int has type int and value "int".
 - The symbol { has the type Ibrace and value "Ibrace".

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Example (Lexical Analysis)

The statement

position = initial + rate * 60;

would be viewed as

$$\mathbf{id}_1 = \mathbf{id}_2 + \mathbf{id}_3 * \mathbf{num};$$

or

id₁ assign id₂ plus id₃ times num semi

by the lexer.

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- There are tools available to assist in the writing of lexical analyzers.
 - lex produces C source code (UNIX).
 - flex produces C source code (gnu).
 - $\bullet \ \mbox{JLex}$ produces Java source code.
 - JFlex produces Java source code.
- We will use JFlex.



The Stages of Compilation

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2 Assignment

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Definition (Syntax analyzer)

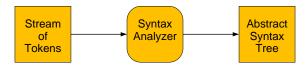
A syntax analyzer, also called a parser, receives a stream of tokens from the lexer and groups them into phrases that match specified grammatical patterns.

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Definition (Abstract syntax tree)

The output of the parser is an abstract syntax tree representing the syntactical structure of the tokens.



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- Grammatical patterns are described by a context-free grammar.
- For example, an assignment statement may be defined as

 $stmt \rightarrow id = expr;$ $expr \rightarrow expr + expr | expr * expr | id | num$

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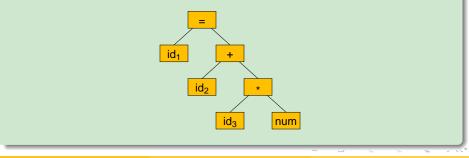
Example

Example (Syntactic Analysis)

• The form

$$id_1 = id_2 + id_3 * num;$$

may be represented by the following syntax tree.



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- There are tools available to assist in the writing of parsers.
 - yacc produces C source code (UNIX).
 - bison produces C source code (gnu).
 - CUP produces Java source code.
- We will use CUP.



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Assignment

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17/33

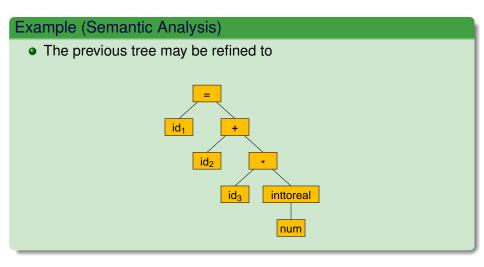
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Definition (Semantic analyzer)

A semantic analyzer traverses the abstract syntax tree, checking that each node is appropriate for its context, i.e., it checks for semantic errors. It outputs a refined abstract syntax tree.

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2 Assignment

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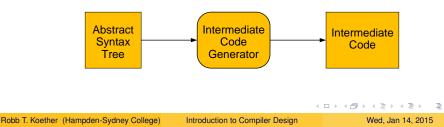
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Definition (Intermediate code)

Intermediate code is code that represents the semantics of a program, but is machine-independent.

Definition (Intermediate code generator)

An intermediate code generator receives the abstract syntax tree and outputs intermediate code that semantically corresponds to the abstract syntax tree.

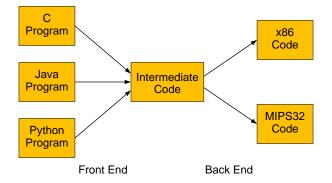


- This stage marks the boundary between the front end and the back end.
- The front end is language-specific and machine-independent.
- The back end is machine-specific and language-independent.

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Intermediate Code



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Example (Intermediate Code Generation)

 The tree in our example may be expressed in intermediate code as

```
temp1 = inttoreal(60)
temp2 = id3 * temp1
temp3 = id2 + temp2
id1 = temp3
```



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2 Assignment

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Definition (Optimizer)

An optimizer reviews the code, looking for ways to reduce the number of operations and the memory requirements.

- A program may be optimized for speed or for size.
- Typically there is a trade-off between speed and size.



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Example (Optimization)

The intermediate code in this example may be optimized as

temp1 = id3 * 60.0
id1 = id2 + temp1

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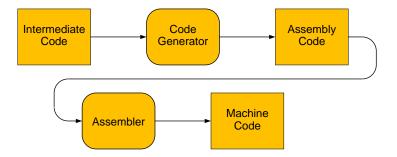
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- The code generator receives the (optimized) intermediate code.
- It produces either
 - Machine code for a specific machine, or
 - Assembly code for a specific machine and assembler.
- If it produces assembly code, then an assembler is used to produce the machine code.

Machine Code Generation



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• The intermediate code may be translated into the assembly code

movf id3,R2
mulf #60.0,R2
movf id2,R1
addf R2,R1
movf R1,id1

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Assignment

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Assignment

• Read Chapters 1 and 2.

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